

# Stability for Take-Away Games

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- $\alpha \geq 1$
- Two players alternate moves
- First turn: take at least one stone, but not all of them
- Subsequent turns: take up to  $\alpha$  times as many stones as last player took
- Goal: Take the last stone. (More precisely, if it's your turn and you can't move, you lose.)



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- $\alpha = 1$ :  $\mathcal{P}$  positions are powers of 2, together with 0
- $\alpha = 2$ :  $\mathcal{P}$  positions are Fibonacci numbers

## $\mathcal{P}$ positions for $\alpha$ -TAG

### Theorem (Schwenk)

The  $\mathcal{P}$  positions satisfy the following recurrence:

$$P_{n+1} = P_n + P_m,$$

where  $m$  is the unique integer such that  $\alpha P_{m-1} < P_n \leq \alpha P_m$ .

Proof relies on a generalization of Zeckendorf's theorem.

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Zeckendorf representation is constructed greedily, by choosing the largest Fibonacci number possible from the remainder.

Example

$$40 = 34 + 5 + 1.$$

# Generalization of Zeckendorf's Theorem

## Theorem (Generalized Zeckendorf's Theorem)

Let  $\alpha \geq 1$ , and let  $P_n$  be a sequence defined by the recurrence  $P_{n+1} = P_n + P_m$ , where  $m$  is the unique integer such that  $\alpha P_{m-1} < P_n \leq \alpha P_m$ , with initial conditions  $P_0 = 0$  and  $P_1 = 1$ . Then every positive integer  $n$  can be expressed uniquely in the form  $n = P_{i_1} + P_{i_2} + \dots + P_{i_k}$ , where  $\alpha P_{i_j} < P_{i_{j+1}}$ .

Construction here is also greedy: take the largest  $P_n$  you can.

## Winning strategy for $\alpha$ -TAG

Suppose there are  $n$  stones. Write  $n = P_{i_1} + P_{i_2} + \cdots + P_{i_k}$  as in the generalized Zeckendorf theorem. Then, on every move, remove the smallest generalized Zeckendorf part  $P_{i_1}$ . The next player will not be able to do so.



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Sample game play with good play by first player:

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# Stability Theorem

## Theorem

*For every  $\alpha > 1$ , there exists a half-open interval  $I_\alpha = [a_\alpha, b_\alpha)$  containing  $\alpha$  such that for all  $\beta \in I_\alpha$ , the  $\mathcal{P}$  positions of  $\alpha$ -TAG are the same as those of  $\beta$ -TAG.*



# Eventual recurrence

## Theorem

*For any  $\alpha \geq 1$ , there exist integers  $k, N \geq 0$  such that for all  $n \geq N$ ,  $P_{n+1} = P_n + P_{n-k}$ .*

However, first few terms might not satisfy this recurrence.

# Stable intervals

Range	Eventual recurrence	Initial conditions
$1 \leq \alpha < 2$	$P_{n+1} = P_n + P_n$	0,1
$2 \leq \alpha < \frac{5}{2}$	$P_{n+1} = P_n + P_{n-1}$	0,1,2
$\frac{5}{2} \leq \alpha < 3$	$P_{n+1} = P_n + P_{n-2}$	0,1,2,3,5
$3 \leq \alpha < \frac{7}{2}$	$P_{n+1} = P_n + P_{n-3}$	0,1,2,3,4,6
$\frac{7}{2} \leq \alpha < \frac{11}{3}$	$P_{n+1} = P_n + P_{n-4}$	0,1,2,3,4,6,8,11,15,21
$\frac{11}{3} \leq \alpha < \frac{43}{11}$	$P_{n+1} = P_n + P_{n-4}$	0,1,2,3,4,6,8,11
$\frac{43}{11} \leq \alpha < 4$	$P_{n+1} = P_n + P_{n-5}$	0,1,2,3,4,6,8,11,14,18,24,32,43
$4 \leq \alpha < \frac{13}{3}$	$P_{n+1} = P_n + P_{n-5}$	0,1,2,3,4,5,7,9,12
$\frac{13}{3} \leq \alpha < \frac{31}{7}$	$P_{n+1} = P_n + P_{n-6}$	0,1,2,3,4,5,7,9,12,15,19,24,31,40,52
$\frac{31}{7} \leq \alpha < \frac{9}{2}$	$P_{n+1} = P_n + P_{n-6}$	0,1,2,3,4,5,7,9,12,15,19,24,31
$\frac{9}{2} \leq \alpha < \frac{14}{3}$	$P_{n+1} = P_n + P_{n-6}$	0,1,2,3,4,5,7,9,11,14,18

# Cutoffs

## Definition

A *cutoff* is some  $\alpha > 1$  such that for every  $\beta < \alpha$ , the sequence of  $\mathcal{P}$ -positions for  $\beta$ -TAG differs from the sequence of  $\mathcal{P}$ -positions for  $\alpha$ -TAG.

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The first few cutoffs are  $2, \frac{5}{2}, 3, \frac{7}{2}, \frac{11}{3}, \frac{43}{11}, 4, \frac{13}{3}, \frac{31}{7}, \frac{9}{2}, \frac{14}{3}$ .

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First  $\mathcal{P}$  positions are 0, 1, 2, 3, 4, 6, 8, 11, 15, 21, 29, 40, 55, ...

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For example, we form 21 as  $15 + 6$ , because  $\alpha \cdot 4 < 15 \leq \alpha \cdot 6$  with  $\alpha = 3$ .



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If we increase  $\alpha$  to  $\frac{15}{4}$ , then the left inequality fails. Thus there is a *potential* cutoff at  $\frac{15}{4}$ , and *definitely* some cutoff in  $(3, \frac{15}{4}]$ .

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Similarly, we compare other terms, to find that the next cutoff is

$$\min \left\{ \frac{4}{1}, \frac{6}{1}, \frac{8}{2}, \frac{11}{3}, \frac{15}{4}, \frac{21}{6}, \frac{29}{8}, \frac{40}{11}, \dots \right\} = \frac{21}{6} = \frac{7}{2}.$$

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## Conjecture

*For every positive integer  $d$ , there exists an integer  $N$  such that for all  $k > N$ ,  $\frac{k}{d}$  is a cutoff.*

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$C(x)$ : number of cutoffs  $\leq x$ . Appears that  $C(x) = \rho x^2 + o(x^2)$  for some  $\rho$ . Probably  $\rho \approx 2$ .



Thank you

Thank you for your attention!